

# Identifying codes in graphs of given maximum degree

Florent Foucaud (LaBRI, Bordeaux, France)

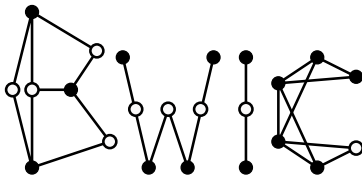
joint works with:

Ralf Klasing, Adrian Kosowski, André Raspaud (2012)

Eleonora Guerrini, Matjaž Kovše, Aline Parreau, Reza Naserasr, Petru Valicov (2011)

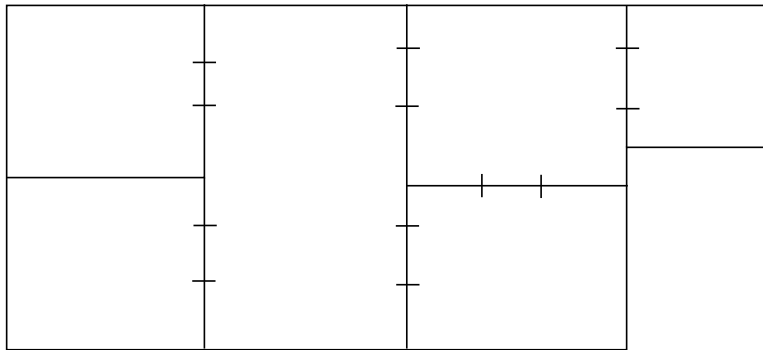
Guillem Perarnau (2012)

Sylvain Gravier, Aline Parreau, Reza Naserasr, Petru Valicov (2011+)

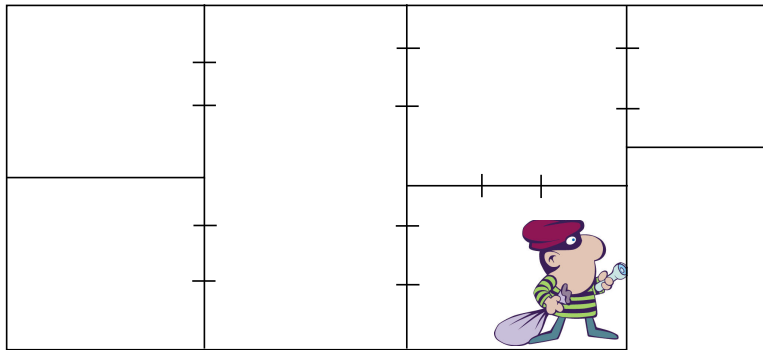


Bordeaux Workshop on Identifying Codes  
November 21-25, 2011

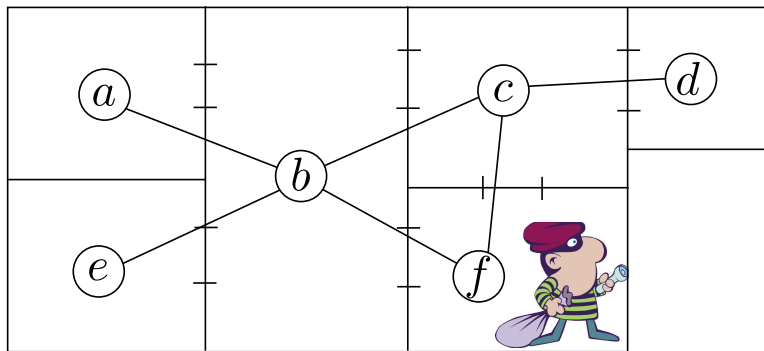
## Locating a burglar in a museum



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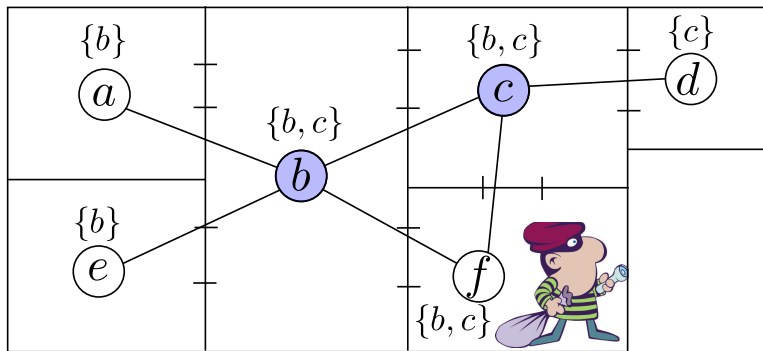


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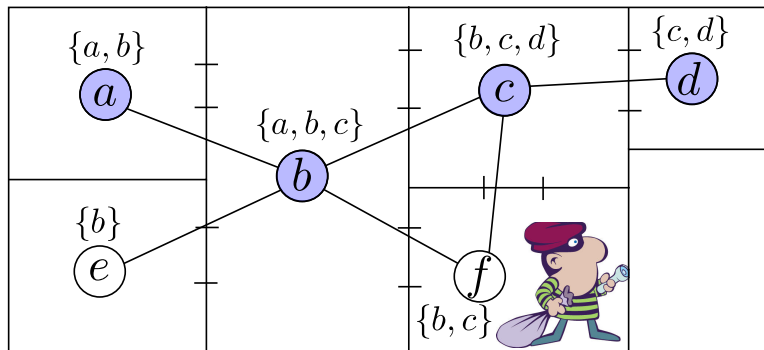
Graph  $G = (V, E)$ .  $V$ : vertices (rooms),  $E \subseteq V \times V$ : edges (doors)

## Locating a burglar in a museum



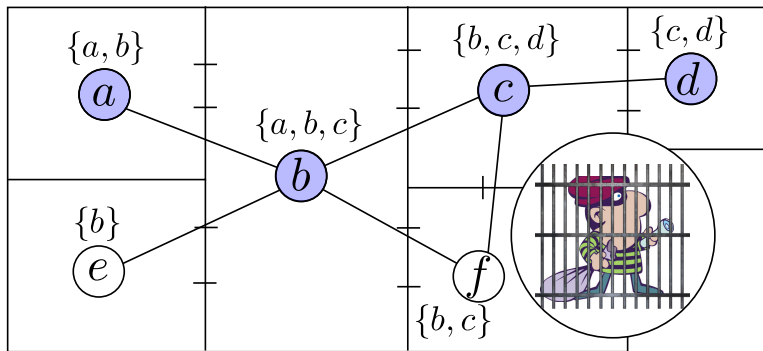
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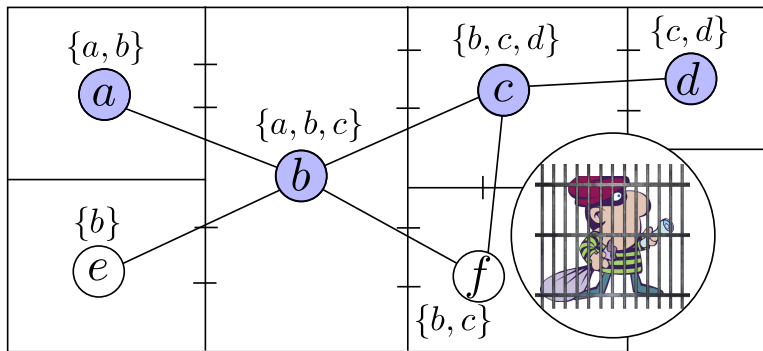
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# Locating a burglar in a museum



How many **detectors** do we need?



Let  $N[u]$  be the set of vertices  $v$  s.t.  $d(u, v) \leq 1$

**Definition** - Identifying code of  $G$  (Karpovsky, Chakrabarty, Levitin, 1998)

Subset  $C$  of  $V$  such that:

- $C$  is a **dominating set** in  $G$ :  $\forall u \in V, N[u] \cap C \neq \emptyset$ , and
- $C$  is a **separating code** in  $G$ :  $\forall u \neq v$  of  $V, N[u] \cap C \neq N[v] \cap C$

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**Notation** - Identifying code number

$\gamma^{\text{ID}}(G)$ : minimum cardinality of an identifying code of  $G$

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**Proposition**

$C$  is an identifying code IFF:

- $C$  is a **dominating set** in  $G$
- $\forall u \neq v$  of  $V$  **with**  $d_G(u, v) \leq 2, (N[u] \Delta N[v]) \cap C \neq \emptyset$

**Theorem** (lower bound: Karpovsky, Chakrabarty, Levitin, 1998  
upper bound: Bertrand, 2005 / Gravier, Moncel, 2007 / Skaggs, 2007)

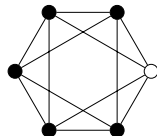
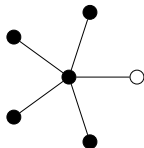
Let  $G$  be an identifiable graph on  $n$  vertices with at least one edge, then

$$\lceil \log_2(n+1) \rceil \leq \gamma^{\text{ID}}(G) \leq n-1$$

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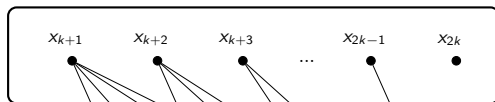
# A class of graphs called $\mathcal{A}$

## Definition - Graph $A_k$

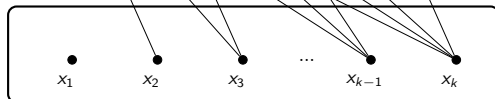
$$V(A_k) = \{x_1, \dots, x_{2k}\}.$$

$x_i$  connected to  $x_j$  iff  $|j - i| \leq k - 1$

Note:  $A_1 = \overline{K_2}$ ; for  $k \geq 2$ ,  $A_k = P_{2k}^{k-1}$



*Clique on  $\{x_{k+1}, \dots, x_{2k}\}$*



*Clique on  $\{x_1, \dots, x_k\}$*

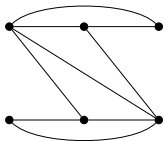
# A class of graphs called $\mathcal{A}$ - examples



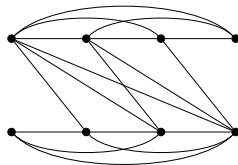
$$A_1 = \overline{K_2}$$



$$A_2 = P_4$$



$$A_3 = P_6^2$$



$$A_4 = P_8^3$$



## Definition - Join and its closure

$(\mathcal{A}, \bowtie)$ : closure of graphs of  $\mathcal{A}$  with respect to  $\bowtie$  (complete join).

## Theorem (F., Guerrini, Kovše, Naserasr, Parreau, Valicov, 2011)

Let  $G$  be an identifiable graph on  $n$  vertices. Then:

$$\gamma^{\text{ID}}(G) = n - 1 \Leftrightarrow G \in \{K_{1,n-1}\} \cup (\mathcal{A}, \bowtie) \cup (\mathcal{A}, \bowtie) \bowtie K_1 \text{ and } G \neq \overline{K_2}.$$

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## Observation

All these graphs have maximum degree  $n - 1$  or  $n - 2$ !

**Theorem** (Karpovsky, Chakrabarty, Levitin, 1998)

Let  $G$  be an identifiable graph with maximum degree  $\Delta$  and  $n$  vertices, then

$$\frac{2n}{\Delta+2} \leq \gamma^{\text{ID}}(G)$$

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**Conjecture** (F., Klasing, Kosowski, Raspaud, 2009)

Let  $G$  be a connected nontrivial identifiable graph on  $n$  vertices and of maximum degree  $\Delta$ . Then:

$$\gamma^{\text{ID}}(G) \leq n - \frac{n}{\Delta} + c \text{ (for some constant } c\text{).}$$

The conjecture is true for  $\Delta = 2$  (with  $c = 3/2$ ).

**Conjecture** (F., Klasing, Kosowski, Raspaud, 2009)

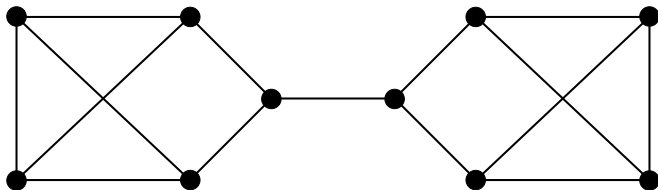
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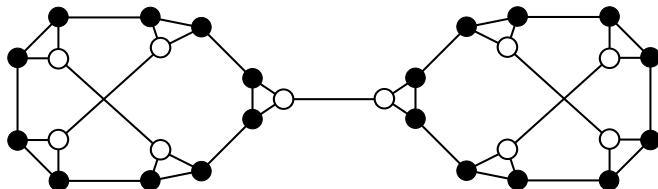
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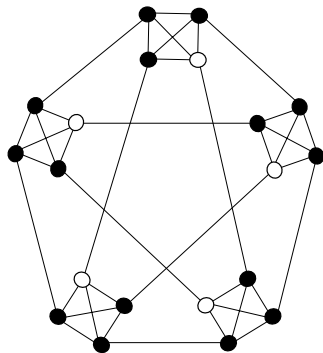
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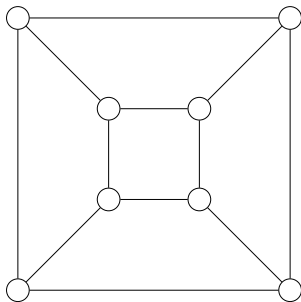




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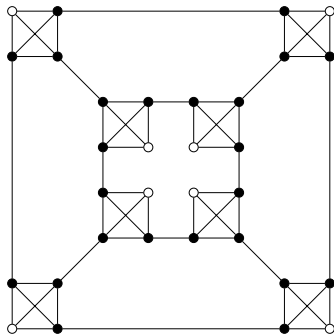
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Also: Sierpiński graphs

(see A. Parreau, S. Gravier, M. Kovše, M. Mollard and J. Moncel, 2011+)

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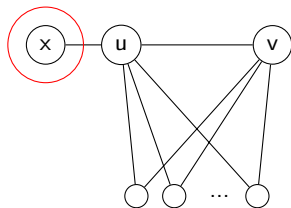
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## Question

Can we prove that  $\gamma^{\text{ID}}(G) \leq n - \frac{n}{\Theta(\Delta)}$ ?

$u, v$  such that  $N[v] \ominus N[u] = \{x\}$

Then  $x \in C$ , forced by  $uv$ .



Note: if  $G$  regular, no forced vertices.

**Theorem** (F., Guerrini, Kovse, Naserasr, Parreau, Valicov, 2011)

Let  $G$  be a connected identifiable graph of maximum degree  $\Delta$ . Then

$$\gamma^{\text{ID}}(G) \leq n - \frac{n}{\Theta(\Delta^5)}$$

If  $G$  is  $\Delta$ -regular,  $\gamma^{\text{ID}}(G) \leq n - \frac{n}{\Theta(\Delta^3)}$

**Proof idea:**

**Proposition**

Let  $I$  be a distance 4-independent set of  $G$ . If for all  $x \in I$ ,  $x$  is **not forced**,  $V - I$  is also an identifying code.

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For each vertex  $x$  of  $G$ , there exists a **non forced vertex**  $y$  in  $N[x]$ .

Take a (maximal) 6-independent set  $I$ . Find the set  $I'$  "good vertices" which are not forced:  $|I| = |I'|$ .  $V - I'$  is an identifying code.

For regular graphs, there are no forced vertices: a 4-IS is enough.



## Theorem (F., Klasing, Kosowski, Raspaud, 2009)

Let  $G$  be a connected identifiable **triangle-free** graph on  $n$  vertices and of maximum degree  $\Delta$ . Then

$$\gamma^{\text{ID}}(G) \leq n - \frac{n}{(1 + \frac{3}{\ln \Delta - 1})\Delta} = n - \frac{n}{(1 + o_{\Delta}(1))\Delta}$$

### Proof idea:

Let  $X$  be the set of vertices having at least some **false twin** (false twins:  $u \not\sim v$  and  $N(u) = N(v)$ ).

- If  $X$  is large, at least  $\frac{|X|}{\Delta}$  vertices can be out of a code and we are done
- Otherwise, build a maximal independent set  $S$  with  $|S| > \frac{\ln \Delta}{\Delta} n$  (using J. Shearer's bound)
- Locally modify  $S$  to get  $S'$ , not too small:  $|S'| \geq |S|/3$
- $V \setminus S'$  is an identifying code

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## Theorem (F., Klasing, Kosowski, Raspaud, 2009)

In fact: let  $G$  be a connected identifiable **triangle-free** graph on  $n$  vertices and of maximum degree  $\Delta$  s.t. for all subgraphs  $H$ ,  $\alpha(H) \geq f(\Delta)n_H$ . Then

$$\gamma^{\text{ID}}(G) \leq n - \frac{n}{\Delta + \frac{3}{f(\Delta)}}$$

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## Corollary

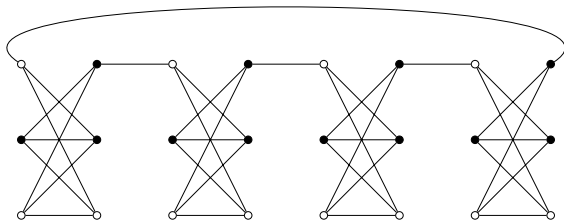
$G$   $k$ -colourable:  $\gamma^{\text{ID}}(G) \leq n - \frac{n}{\Delta + 3k}$ .

$\Rightarrow$  Bipartite:  $\gamma^{\text{ID}}(G) \leq n - \frac{n}{\Delta + 6}$

$\Rightarrow$  Planar triangle-free:  $\gamma^{\text{ID}}(G) \leq n - \frac{n}{\Delta + 9}$

## Triangle-free graphs - examples

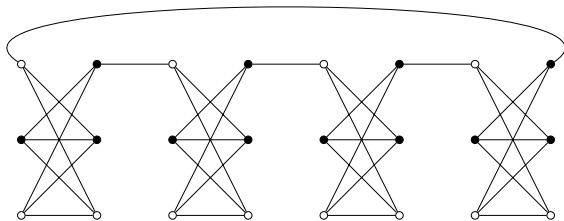
Complete  $(\Delta - 1)$ -ary tree, caterpillar: roughly,  $\gamma^{\text{ID}}(G) = n - \frac{n}{\Delta - 1}$



$$\gamma^{\text{ID}}(G) = n - \frac{n}{2\Delta/3}$$

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$$\gamma^{\text{ID}}(G) = n - \frac{n}{2\Delta/3}$$

### Question

What about triangle-free graphs without false twins?

## Notation

Let  $NF(G)$  be the proportion of **non** forced vertices of  $G$

$$NF(G) = \frac{\text{\#non-forced vertices in } G}{\text{\#vertices in } G}$$

## Theorem (F., Perarnau, 2011)

For each identifiable graph  $G$  on  $n$  vertices having maximum degree  $\Delta \geq 3$  and no isolated vertices,

$$\gamma^{\text{ID}}(G) \leq n - \frac{n \cdot NF(G)^2}{103\Delta}$$

### Proof idea:

- Take all forced vertices (set  $F$ ) into the code.
- From  $V \setminus F$ , select each vertex with probability  $p_S = \frac{1}{k \cdot \Delta}$  ( $k$  constant) to belong to a set  $S$ . We want  $C = V \setminus S$ .
- Use Lovász' Local Lemma to show that  $\Pr(C \text{ is a code}) > f(k, n, \Delta) > 0$
- Use the Chernoff bound to show that  $\Pr(C \text{ is too small}) < f(k, n, \Delta)$

### Proposition

$$\frac{1}{\Delta+1} \leq NF(G) \leq 1$$

**Proof:**

### Lemma (Bertrand, Hudry, 2001)

Let  $G$  be an identifiable graph having no isolated vertices. Let  $x$  be a vertex of  $G$ . There exists a **non forced vertex**  $y$  in  $N[x]$ .

$\Rightarrow$  The set  $S$  of non-forced vertices forms a dominating set. Hence  $|S| \geq \frac{n}{\Delta+1}$ .

## Proposition

Let  $G$  be a graph of **clique number** at most  $k$ . There exists a function  $\rho$  such that:

$$\frac{1}{\rho(k)} \leq NF(G) \leq 1$$

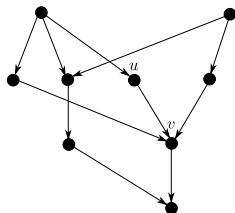


## Proposition

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- Define graph  $\vec{H}(G)$
- Max. degree of  $\vec{H}(G)$ :  $2k - 3$
- Longest directed chain of  $\vec{H}(G)$ :  $k - 1$
- Each component has a non-forced vertex
- $\Rightarrow \rho(k) \leq \sum_{i=0}^{k-2} (2k - 3)^i$



$$u \rightarrow v \text{ if } N[v] = N[u] \cup \{x\}$$

**Theorem** (F., Perarnau, 2011)

For each identifiable graph  $G$  on  $n$  vertices having maximum degree  $\Delta \geq 3$  and no isolated vertices,

$$\gamma^{\text{ID}}(G) \leq n - \frac{n \cdot NF(G)^2}{103\Delta}$$

**Corollary**

- In general,  $NF(G) \geq \frac{1}{\Delta+1}$  and  $\gamma^{\text{ID}}(G) \leq n - \frac{n}{\Theta(\Delta^3)}$
- If  $G$  is  $\Delta$ -regular,  $NF(G) = 1$  and  $\gamma^{\text{ID}}(G) \leq n - \frac{n}{103\Delta} = n - \frac{n}{\Theta(\Delta)}$
- If  $G$  has clique number bounded by  $k$ ,  $NF(G) \geq \frac{1}{\rho(k)}$  and  $\gamma^{\text{ID}}(G) \leq n - \frac{n}{103 \cdot (\rho(k))^2 \cdot \Delta} = n - \frac{n}{\Theta(\Delta)}$

Note: for  $k = 2, 3, 4, 5$ :  $103 \cdot (\rho(k))^2 = 103, 1.360, 81.685, 13.600.000$

The conjecture holds for some large subclass of line graphs:

**Theorem** (F., Gravier, Naserasr, Parreau, Valicov, 2011)

Let  $G$  be an edge-identifiable graph with a minimal edge-identifying code  $C_E$ . Then  $G[C_E]$  is 2-degenerate.

**Corollary**

If  $G$  edge-identifiable,  $\gamma^{\text{ID}}(\mathcal{L}(G)) \leq 2|V(G)| - 3$ .

**Corollary**

If  $G$  is an edge-identifiable graph **with average degree**  $\bar{d}(G) \geq 5$ , then  $\gamma^{\text{ID}}(\mathcal{L}(G)) \leq n - \frac{n}{\Delta(\mathcal{L}(G))}$  where  $n = |V(\mathcal{L}(G))|$ .

## Conjecture (F., Klasing, Kosowski, Raspaud, 2009)

Let  $G$  be a connected nontrivial identifiable graph on  $n$  vertices and of maximum degree  $\Delta$ . Then:

$$\gamma^{\text{ID}}(G) \leq n - \frac{n}{\Delta} + c \text{ (for some constant } c\text{).}$$

- Can we reduce the constants?
- Can we improve the bound  $n - \frac{n}{\Theta(\Delta^3)}$ ?
- What about  $\Delta = 3$ ?
- What about trees (having a look at David Auger's algorithm)?
- What about claw-free graphs?  $n - \frac{n}{\Theta(\Delta^2)}$  seems to hold by directly using similar arguments than for triangle-free graphs.
- Other related parameters?